# Applied Combinatorics by Fred S. Roberts and Barry Tesman

# Answers to Selected Exercises<sup>1</sup>

## Chapter 11

Section 11.1.

**1(a)**. One example: assign labels a = 1, e = 2, j = 3, k = 4, f = 5, g = 6, h = 7, l = 8, m = 9, i = 10, d = 11, c = 12, b = 13 and mark the edges  $\{a, e\}, \{e, j\}, \{j, k\}, \{k, f\}, \{f, g\}, \{g, h\}, \{h, l\}, \{l, m\}, \{m, i\}, \{i, d\}, \{d, c\}, \{c, b\};$ 

2. not connected;

**3(a)**. connected;

4(a). one example: use the marked edges in answer to 1(a);

**6**. it is a spanning forest;

8(a). no;

10. yes.

Section 11.2.

1(a). none;

 $1(f). \{c, e\};$ 

2(a). orientation based on answer to Exercise 1(a), Section 11.1 orients 1 to 2 to 3 to ... to 13 and all other edges from higher number to lower number;

**7**.  $V = \{x, a, b, c\}, E = \{\{x, a\}, \{x, b\}, \{x, c\}\};$ 

8(a). none;

8(f). c, e;

16. for digraph (a) and measure (1): essentially the only orientations are: (i) which uses arcs (a, b), (b, c), (c, f), (f, e), (e, d), (d, a), and (b, e), or (ii) which uses arcs (a, b), (b, e), (e, d), (d, a), (e, f), (f, c), and (c, b); both are equally efficient; for digraph (a) and measure (3): orientation (ii) above is best; for digraph (a) and

<sup>&</sup>lt;sup>1</sup>More solutions to come. Comments/Corrections would be appreciated and should be sent to: Barry Tesman (tesman@dickinson.edu) or Fred Roberts (froberts@dimacs.rutgers.edu).

measure (4): orientation (ii) above is best; for digraph (a) and measure (6): orientation (ii) above is best;

**25**.  $D_1$ : category 3,  $D_4$ : category 2;

**26.** if  $V = \{a, b, c, d\}$  and  $A = \{(a, b), (b, c), (c, d), (d, a)\}$ , any arc is (3, 2);

**27**. in previous example, any vertex is (3, 2);

28(a). 1.

Section 11.3.

**2**.  $G_1$ : none;  $G_2$ :  $a, c, e, b, f, e, i, j, k, l, i, k, h, g, f, d, a; G_3$ : none;  $G_4$ : f, e, c, d, e, b, a, b, a, d, g, f;

**3**.  $G_1$ ;

**4**.  $D_1$ :  $a, b, d, c, d, f, e, c, a; D_2$ : none;  $D_3$ : none;  $D_4$ : none;  $D_5$ : none;

**5**.  $D_2: c, a, b, d, f, e, c, d; D_5: b, a, e, f, g, b, c, d, e;$ 

10(a). yes;

10(b). yes;

10(c). no;

10(d). yes;

 $12(a). D_1: 2;$ 

**12(b)**.  $D_2: 2; D_5: 2;$ 

**13**. 12;

**16**. no: consider  $D_2$  of Figure 11.25.

Section 11.4.

- $\mathbf{3(a)}$ . add edge  $\{c, f\}$ ;
- **3(b)**. add edges  $\{e, g\}$  and  $\{g, h\}$ ;
- $\mathbf{3(c)}$ . add edges  $\{a, d\}$  and  $\{b, c\}$ ;
- $\mathbf{3}(\mathbf{d})$ . add edges  $\{a, d\}$  and  $\{b, c\}$ ;

6(a). CAAGCUGGUC;

**9(a)**. yes:  $A_1A_1A_2A_2A_2A_3A_3A_3A_3A_3A_3A_2A_1$ ;

**17(a)**. say *B* is an interior extended base of a U, C (G) fragment; then both *B* and the preceeding extended base end in G (U, C), so *B* is on the second list;

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### Answers to Selected Exercises

**18**. ends in A;

20(b). if there is a second abnormal fragment, it is B alone.

#### Section 11.5.

**1(a)**. a, b, d, f, e, c, a;**1(b)**. i, a, b, c, d, e, f, g, h, i;**1(c)**. *a*, *b*, *c*, *d*, *h*, *g*, *f*, *j*, *k*, *l*, *p*, *o*, *n*, *m*, *i*, *e*, *a*; 2(a). e, c, d, a, b;**2(b)**. *a*, *b*, *c*, *d*, *f*, *e*; **2(c)**. a, b, c, d, e, j, g, i, f, h;**3(a)**. *a*, *b*, *c*, *d*, *e*, *f*, *a*; **3(b)**. *a*, *b*, *c*, *e*, *d*, *a*; **3(c)**. *c*, *a*, *d*, *b*, *e*, *f*, *c*; **4(a)**. *a*, *d*, *b*, *c*; **4(b)**. *a*, *b*, *d*, *c*, *e*, *f*; **4(c)**. *a*, *c*, *d*, *b*, *e*; 6. no; **7(a)**.  $Z_4$ ; 7(b). K<sub>4</sub>; 7(c). the graph in Figure 11.50; 7(d).  $K_{1,4};$ 8(a).  $K_{1,3}$ ; 8(b). K<sub>4</sub>; 9(a). yes; 10(a). for (a) of Figure 11.45: complete graph; **10(b)**. for (a) of Figure 11.45: yes; **11**.  $K_{1,3}$ ; **12**.  $Z_5$ ; **13**. for (a) of Figure 11.4: no;

**14**. for (a) of Figure 11.4: no;

**15(a)**. use Theorem 11.8.

Section 11.6.

1(b). yes;

**2**. for (b) of Exercise 1, Section 4.1: no;

**3**. for (a), the labeling is a topological order;

8. SF, B, H, LA, NY or SF, B, LA, NY H, or SF, B, NY H, LA;

10. *i* beats *j* iff i < j;

**11(a)**. a = 1, b = 2, c = 3, d = 4, e = 5, f = 6;

**12**. 3, 1, 2;

**15.** if C is a, b, f, d, c, e, a, then H has edges  $\{b, d\}$  to  $\{e, f\}$  and  $\{a, c\}$  to  $\{e, f\}$  and is 2-colorable;

**20(b)**. (0, 1, 2, 3, 4);

22(a). no;

22(b). no;

**23**. consider the tournament on  $V(D) = \{1, 2, 3, 4\}$  and  $A(D) = \{(1, 2), (2, 4), (3, 1), (3, 2), (4, 1), (4, 3)\};$ 

**28(a)**.  $(0, 1, 2, \ldots, n-1);$ 

**31**. start with any vertex x and find the longest simple path heading into x; this must start at a vertex with no incoming arcs;

**32(a)**.  $\binom{s(u)}{2};$ 

**32(c)**. use part (b) and the fact that  $s(u) \ge 2$  for some vertex u in every tournament of four or more vertices.